

The Hardware & Software Implications of Microservices and How Big Data Can Help

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Executive Summary



- Shift from monoliths to microservices:
 - Modularity, specialization, simplicity, accelerated development
 - Change assumptions about datacenter server design
 - Complicate scheduling and resource management
 - Amplify tail@scale effects

- Revisit architectural design decisions for microservices
- Highlight management challenges of microservices
- Motivate the need for data-driven approaches for systems whose scale & complexity keeps increasing

From Monoliths to Microservices



Motivation

Advantages of microservices:

- Ease & speed of code development & deployment
- Security, error isolation
- PL/framework heterogeneity

Challenges of microservices:

- Change server design assumptions
- \square Complicate resource management \rightarrow dependencies
- Amplify tail-at-scale effects
- More sensitive to performance unpredictability
- No representative end-to-end apps with microservices

An End-to-End Suite for Cloud & IoT Microservices

- □ 4 end-to-end applications using popular open-source microservices → ~30-40 microservices per app
 - Social Network
 - Movie Reviewing/Renting/Streaming
 - E-commerce
 - Drone control service

Programming languages and frameworks:

- node.js, Python, C/C++, Java/Javascript, Scala, PHP, and Go
- Nginx, memcached, MongoDB, CockroachDB, Mahout, Xapian
- Apache Thrift RPC, RESTful APIs
- Docker containers
- Lightweight RPC-level distributed tracing



Movie Streaming

- Browse movie info (movie plot, photos, videos, cast, stats, etc.)
- ML widgets:
 - Recommender for movies to watch
 - Recommender for ads
- User authentication/Payment
- Search:
 - Xapian: search movie DB
- Analytics:
 - Mahout: user analytics based on input stored in HDFS
 - Spark MLlib: in-memory ML analytics

[CAL'18]

MongoDB





Memcached

Big vs. small servers:

- Power management using RAPL
- More pressure on single-thread performance, low tail latency





Big vs. small servers:

- Power management using RAPL
- More pressure on single-thread performance, low tail latency
- Low-power SoCs, e.g., Cavium ThunderX2
- Similar latency, but earlier saturation



Computation:Communication ratio:

- Monolithic service \rightarrow 70:30 @ high load
- Microservices \rightarrow 50:50 @ high load





Computation:Communication ratio:

- Monolithic service → 70:30 @ high load
- Microservices → 50:50 @ high load
- **\square** RPC/REST acceleration \rightarrow NIC offloads, FPGAs



L1-i cache pressure:

- \square Monoliths \rightarrow Large code footprints \rightarrow L1i thrashing
- $\square \text{ Microservices} \rightarrow \text{Small footprint/microservice}$
 - Assuming dedicated cores

End-to-End Latency Breakdown



Post-rightsizing (resource ratios to avoid glaring bottlenecks)

- Bottlenecks shift with load
- Need online, dynamic decisions

Resource Management Implications



□ Challenges of microservices:

- Change server design assumptions
- Dependencies complicate resource management

Dependencies & Backpressure



Dependencies & Backpressure



- Traditional techniques like autoscale may help/penalize the wrong microservice
- \Box Dependencies change at runtime \rightarrow difficult to infer impact

Determine Per-Tier QoS



Queueing network simulation

Complex microservices graphs, blocking, cyclic dependencies, etc.

Power Management for Microservices

Two types of latency slack:

Microservices off the critical path

Microservices on the critical path with relaxed QoS



Scalability Challenges



Tail at Scale Effects

□ Microservices add an extra dimension to tail at scale effects

- A single slow microservice affects end-to-end latency
- Much more pressure on performance predictability & availability
- Monitoring at the edge

 Determining per-tier QoS for 10000s of microservices is intractable

Scalable data-driven approach

Need for online performance debugging

Proactive Performance Debugging

- - Finding the culprit of a QoS violation is difficult
 - Post-QoS violation, returning to nominal operation is hard
- Anticipating QoS violations & identifying culprits
- Seer: Data-driven Performance Debugging for Microservices
 - Combines lightweight RPC-level distributed tracing with hardware monitoring
 - Leverages scalable deep learning to signal QoS violations with enough slack to apply corrective action

Queue OCPU Mem Net Disk Performance Implications

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Queue OCPU Mem Net Disk Performance Implications

Seer: Data-Driven Performance Debugging

[HotCloud'18]

Leverage the massive amount of traces collected over time

- Apply online, practical data mining techniques that identify the culprit of an *upcoming* QoS violation
- 2. Use per-server hardware monitoring to determine the cause of the QoS violation
- 3. Take corrective action to prevent the QoS violation from occurring

□ Need to predict 100s of msec – a few sec in the future



Deep Learning to the Rescue

□ Why?

- Architecture-agnostic
- Adjusts to changes in dependencies over time
- High accuracy, good scalability
- Inference within the required window







Training once: slow (hours - days)

- Across load levels, load distributions, request types
- Distributed queue traces, annotated with QoS violations
- Weight/bias inference with SGD
- Retraining in the background
- Inference continuously: streaming trace data







Challenges:

- In large clusters inference too slow to prevent QoS violations
- Offload on TPUs, 10-100x improvement; 10ms for 90th %ile inference
- Fast enough for most corrective actions to take effect (net bw partitioning, RAPL, cache partitioning, scale-up/out, etc.)

Experimental Setup

- 40 dedicated servers
- ~1000 single-concerned containers
- Machine utilization 80-85%
- Inject interference to cause
 QoS violation
 - Using microbenchmarks (CPU, cache, memory, network, disk I/O)



Restoring QoS

Identify cause of QoS violation

- Private cluster: performance counters & utilization monitors
- Public cluster: contentious microbenchmarks

Adjust resource allocation

- RAPL (fine-grain DVFS) & scale-up for CPU contention
- Cache partitioning (CAT) for cache contention
- Memory capacity partitioning for memory contention
- Network bandwidth partitioning (HTB) for net contention
- Storage bandwidth partitioning for I/O contention

Queue OCPU Mem ()Net ()Disk

Seer

Demo

Default



Demo

A New Cloud Stack



A New Cloud Stack

